

**EE 193 - Applied Probability and Statistics for Engineers**  
**Department of Electrical and Computer Engineering**  
**Tufts University Fall 2007**  
**Problem Set #3 Solutions**

**Problem 1**

Yates and Goodman problems

• **2.2.3**

(a) Since the sum of  $P_V(v)$  over all  $v$  must equal 1 we have

$$\sum_{v=1}^4 P_V(v) = c [1^2 + 2^2 + 3^2 + 4^2] = 1 \rightarrow 30c = 1 \rightarrow c = \frac{1}{30}$$

(b) This is the probability that  $v$  is a square of an integer. Since  $v$  can assume values from 1 to 4, the only two squares are 1 and 4.

$$P[V \in \{u^2 | u = 1, 2, \dots\}] = P_V(1) + P_V(4) = \frac{1}{30} + \frac{16}{30} = \frac{17}{30}$$

(c)  $P[V \text{ is even}] = P[V = 2 \cup V = 4] = \frac{4}{30} + \frac{16}{30} = \frac{20}{30} = \frac{2}{3}$

(d)  $P[V > 2] = P[V = 3 \cup V = 4] = \frac{25}{30} = \frac{5}{6}$

- **2.2.6** Each call is a Bernoulli random variable with probability  $p$ . Failing to serve a person who calls three times means that person experiences three failures in a row. The probability of this occurring is  $(1 - p)^3$ . Thus, the probability that the person is served is  $1 - (1 - p)^3$ . This must equal 0.95. Hence we require

$$1 - (1 - p)^3 = 0.95 \rightarrow (1 - p)^3 = 0.05 \rightarrow p = 1 - .05^{1/3} \approx 0.63$$

- **2.3.2** A Bernoulli process model is appropriate here where a “success” is the reception of a message on one trial. The probability of a success is  $p$ .

(a) The probability we obtain  $K$  successes in  $n$  trials is a binomial random variable

$$P_K(k) = \binom{n}{k} p^k (1 - p)^{n-k} \quad k = 0, 1, 2, \dots$$

(b) The probability of receiving the message at least once is 1 minus the probability of receiving the message zero times.

$$\begin{aligned} P[\text{At least once}] &= 1 - p_K(0) = 1 - \binom{n}{0} (0.8)^0 (0.2)^n \\ &= 1 - 0.2^n \end{aligned}$$

The requirement we have then is  $1 - 0.2^n = 0.95$  or  $0.2^n = .05$ . Taking the natural log of both sides and dividing gives

$$\frac{\ln 0.05}{\ln 0.2} = 1.86$$

Taking  $n$  to be the first integer larger than 1.86, that is 2, will give the required result.

- **2.3.11** The packet is transmitted once with probability  $p$ . For it to be transmitted twice, we need one failure and then one success for a probability of  $(1-p)p$ . For the packet to be transmitted three times, we need two failures followed by one success. This probability is  $(1-p)^2p$ . As long as the number of tries is less than  $d$ , the probability is geometric. Now, if we have failed at time  $d$  then we stop transmitting. So the probability that we transmit  $d$  times is  $(1-p)^d$ . Thus it would appear that

$$p_T(t) = \begin{cases} p(1-p)^{t-1} & t = 0, 1, 2, \dots, d-1 \\ (1-p)^t & t = d \\ 0 & \text{else} \end{cases}$$

As a sanity check (and as an exercise in algebra), let's check that this sums to 1.

$$\begin{aligned} \sum_{t=1}^d p_T(t) &= p + p(1-p) + p(1-p)^2 + \dots + p(1-p)^{d-1} + (1-p)^d \\ &= \left[ \sum_{t=0}^{d-1} p(1-p)^t \right] + (1-p)^d = \left[ p \frac{1 - (1-p)^d}{1 - (1-p)} \right] + (1-p)^d = 1 \end{aligned}$$

where in simplifying the sum in the second line, I have used the formula

$$\sum_{i=m}^n r^i = \frac{1 - r^{n+1}}{1 - r}$$

taking  $r = 1 - p$ ,  $m = 0$  and  $n = d - 1$ .

- **2.4.5** This is similar to problem 2.3.11. If the first pizza is mushroom, then with probability  $p = 2/3$ , there are zero pizzas sold before the first "success." If the first pizza is not mushroom but the second pizza is, then  $N = 1$  with probability  $(1-p)p$ . More generally, the probability that  $N$  pizzas were sold before the first mushroom is equal to  $(1-p)^N p$ . This works for  $0 \leq N < 99$ . With probability  $(1-p)^{100}$ , there will be 100 pizzas sold before the first mushroom. Thus, the PMF is

$$p_N(n) = \begin{cases} p(1-p)^n & n = 0, 1, 2, \dots, 99 \\ (1-p)^{100} & n = 100 \\ 0 & \text{else} \end{cases}.$$

For  $x < 0$ , the CDF is zero. For integers between 0 and 100 the CDF is

$$F_N(n) = \sum_{i=0}^n p(1-p)^i = p \frac{1 - (1-p)^{n+1}}{1 - (1-p)} = 1 - \left(\frac{1}{3}\right)^{n+1}.$$

Thus using the "floor" notation from section 2.4 of the text, for  $0 \leq x < 100$

$$F_N(x) = 1 - \left(\frac{1}{3}\right)^{\lfloor x \rfloor + 1}.$$

Finally for  $x \geq 100$  the CDF is 1.

• 2.5.2

- (a) From the question, a call will cost 20 cents (since it is voice) with probability 0.6 or it will cost 30 cents (since the call is data) with probability 0.4. Hence

$$p_C(c) = \begin{cases} 0.6 & c = 20 \\ 0.4 & c = 30 \\ 0 & \text{else} \end{cases}$$

- (b) Using the definition of expected value:  $E[C] = \sum_c cp_C(c) = 20 \times 0.6 + 30 \times 0.4 = 24$

- 2.5.9 The tree diagram for this problem is shown in Fig. 1. The top branches are the amount we take home after a win. The bottom branches show the progress of the game when we loose. All edges have probabilities equal to  $1/2$ . As we see, we either walk away with \$64 (regardless of how long it takes to win), or we loose everything and go home with \$0. The probability of loosing everything is  $(\frac{1}{2})^6 \approx 0.016$ . The probability that we walk away with \$64 is  $1 - (\frac{1}{2})^6 \approx 0.984$ . The expected value of the amount we take home is  $64 \times 1 - (0.5)^6 = 62$ . Since this is the same as the amount we started with, this is not a game we would choose to play.

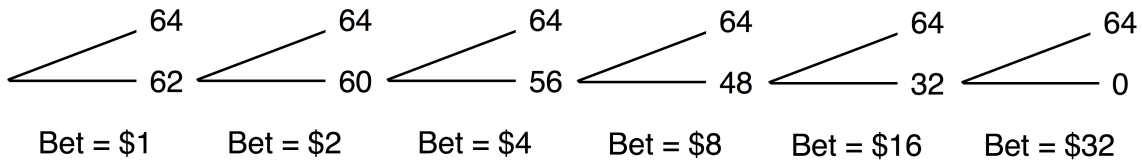


Figure 1: Problem 2.5.9

- 2.6.2 First, from the CDF of  $X$  in Problem 2.4.2, we find that the PDF is

$$p_X(x) = \begin{cases} 0.2 & x = -1 \\ 0.5 & x = 0 \\ 0.3 & x = 1 \\ 0 & \text{else} \end{cases} .$$

Now, since  $V = |X|$ , we have that  $V = 1$  for  $X = -1$  as well as  $X = 1$  and  $V = 0$  when  $X = 0$ . Hence  $p_V(1) = p_X(-1) + p_X(1) = 0.5$  and  $p_V(0) = p_X(0) = 0.5$  so

$$p_V(v) = \begin{cases} 0.5 & v = 0 \\ 0.5 & v = 1 \\ 0.0 & \text{else} \end{cases}$$

and

$$F_V(x) = \begin{cases} 0 & x < 0 \\ 0.5 & 0 \leq x < 1 \\ 1 & x \geq 1 \end{cases} .$$

Finally,  $E[V] = 0 \times 0.5 + 1 \times 0.5 = 0.5$ .

• 2.6.5

- (a) Each transmission is a Bernoulli trial where the probability of error is  $q$ . The number of trials before the first non-error is thus geometric where, rather than waiting for the first “success” (here transmission with error) we are instead thinking about the first failure (no error)

$$p_X(x) = q^{x-1}(1 - q) \quad x = 1, 2, 3, \dots$$

- (b) Let’s start by enumerating what is going on. When  $X = 1$  we are successful with the first attempt and  $T = 1$ . When  $X = 2$ , the sequence of events is transmit-NACK-transmit again so  $T = 3$ . Similarly, when  $X = 3$  we have to go through 2 transmit-NACK cycles and then a fifth, successful transmit so  $T = 5$ . In general then  $T = 2X - 1$  and  $p_T(t) = p_X((t + 1)/2)$  or

$$p_T(t) = q^{(t+1)/2-1}(1 - q) \quad t = 1, 3, 5, \dots$$