

Warmup

Write a boolean equation in canonical form for this truth table:

A	B	C		Y
0	0	0		0
0	0	1		0
0	1	0		0
0	1	1		1
1	0	0		0
1	0	1		1
1	1	0		0
1	1	1		1

$$\bar{A} B C + A \bar{B} C + A B C = Y$$

ES 4: Boolean algebra

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By the end of class today, you should be able to:

- Use boolean algebra to manipulate equations, using theorems 1-12
- Use "bubble pushing" to manipulate the logic of a circuit
- Use the above manipulation techniques to implement a circuit with constraints (e.g., only NAND gates)

Key representations

and how to simplify them

Boolean equation

$$\overline{AB} + C$$

Boolean algebra

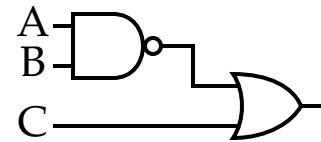
Truth table

A	B	C	Y
0	0	0	1
0	0	1	1
0	1	0	1

Karnaugh maps

(next class)

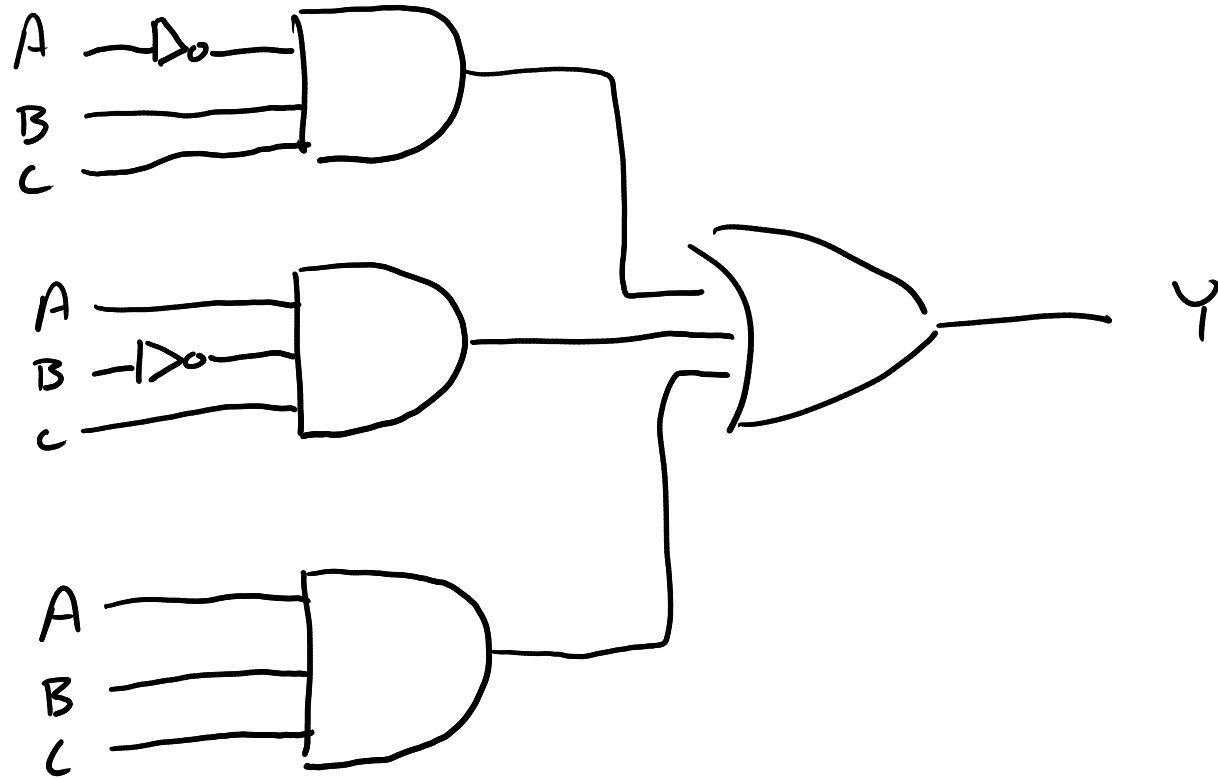
Logic diagram



Bubble pushing

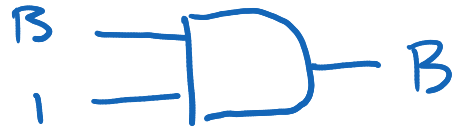
What if we implemented this with gates?

$$\overline{A}BC + A\overline{B}C + ABC$$

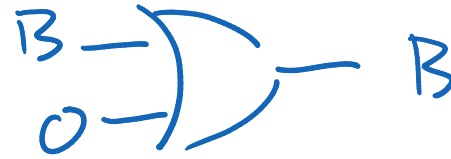


Boolean theorems to the rescue! (??)

$$B \cdot 1 = B$$



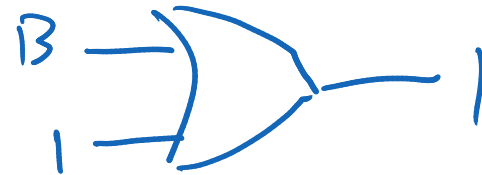
$$B + 0 = B$$



$$B \cdot 0 = 0$$



$$B + 1 = 1$$



More Boolean theorems to the rescue!

$$B \cdot B = \mathcal{B}$$

$$B + B = \mathcal{B}$$

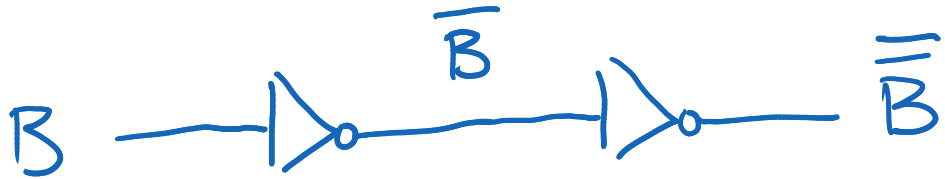
$$B \cdot \overline{B} = 0$$

$$B + \overline{B} = 1$$

Inverting twice

(aka "involution")

$$\overline{\overline{B}} = B$$



Associativity

Works just like "normal" algebra

$$A B C = (A B) C = A (B C)$$

$$A + B + C = (A + B) + C = A + (B + C)$$

Distributivity, what?

AND distributes over **OR**

$$A(B + C) = AB + AC$$

But also: **OR** distributes over **AND**

$$A + (BC) = (A + B)(A + C)$$

Simplifying from canonical form

$$\overline{A}BC + \overline{A}\overline{B}C + A\overline{B}C + ABC$$

$$\overline{A}(BC) + A(\overline{B}C) + A\overline{B}C + ABC$$

$$(\overline{A} + A)(BC) + A\overline{B}C + ABC$$

$$BC + (\overline{B} + B)(AC)$$

$$BC + AC$$

$$(A + B)C$$

Practice

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	1
0	0	1	0	0
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0

A	B	C	D	Y
1	0	0	0	1
1	0	0	1	1
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

$$\bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D + A\bar{B}\bar{C}\bar{D} + A\bar{B}\bar{C}D + A\bar{B}CD$$

$$\bar{B}(\bar{A}\bar{C}\bar{D} + \bar{A}\bar{C}D + A\bar{C}\bar{D} + A\bar{C}D + ACD)$$

$$\bar{B}(\bar{A}\bar{C} + A\bar{C} + ACD + A\bar{C}D)$$

$$\bar{B}((\bar{A} + A)(\bar{C}) + (AD)(\bar{C} + C))$$

$$\bar{B}(\bar{C} + AD)$$

All the theorems

$$T1 \quad B \cdot 1 = B \quad B + 0 = B$$

$$T2 \quad B \cdot 0 = 0 \quad B + 1 = 1$$

$$T3 \quad B \cdot B = B \quad B + B = B$$

$$T4 \quad \overline{\overline{B}} = B$$

$$T5 \quad B \cdot \overline{B} = 0 \quad B + \overline{B} = 1$$

$$T6 \quad B \cdot C = C \cdot B \quad B + C = C + B$$

$$T7 \quad (B \cdot C) \cdot D = B \cdot (C \cdot D)$$

$$(B + C) + D = B + (C + D)$$

$$T8 \quad (B \cdot C) + (B \cdot D) = B \cdot (C + D)$$

$$(B + C) \cdot (B + D) = B + (C \cdot D)$$

$$T9 \quad B \cdot (B + C) = B$$

$$B + (B \cdot C) = B$$

$$T10 \quad (B \cdot C) + (B \cdot \overline{C}) = B$$

$$(B + C) \cdot (B + \overline{C}) = B$$

$$T11 \quad (B \cdot C) + (\overline{B} \cdot D) + \cancel{(C \cdot D)}$$

$$(B + C) \cdot (\overline{B} + D) \cdot \cancel{(C + D)}$$

$$T12 \quad \overline{B_0 \cdot B_1 \cdot B_2 \dots} = \overline{B_0} + \overline{B_1} + \overline{B_2} \dots \quad \overline{B_0 + B_1 + B_2 \dots} = \overline{B_0} \cdot \overline{B_1} \cdot \overline{B_2} \dots$$

What do I need to remember?

The axioms

Theorems 1-8, and DeMorgan's theorem (12)

What do I need to ~~remember?~~

be able to use?

The axioms

Theorems 1-8, and DeMorgan's theorem (12)

Intermission: lab notes

With lab notes, we're trying to push you toward *professional practice*

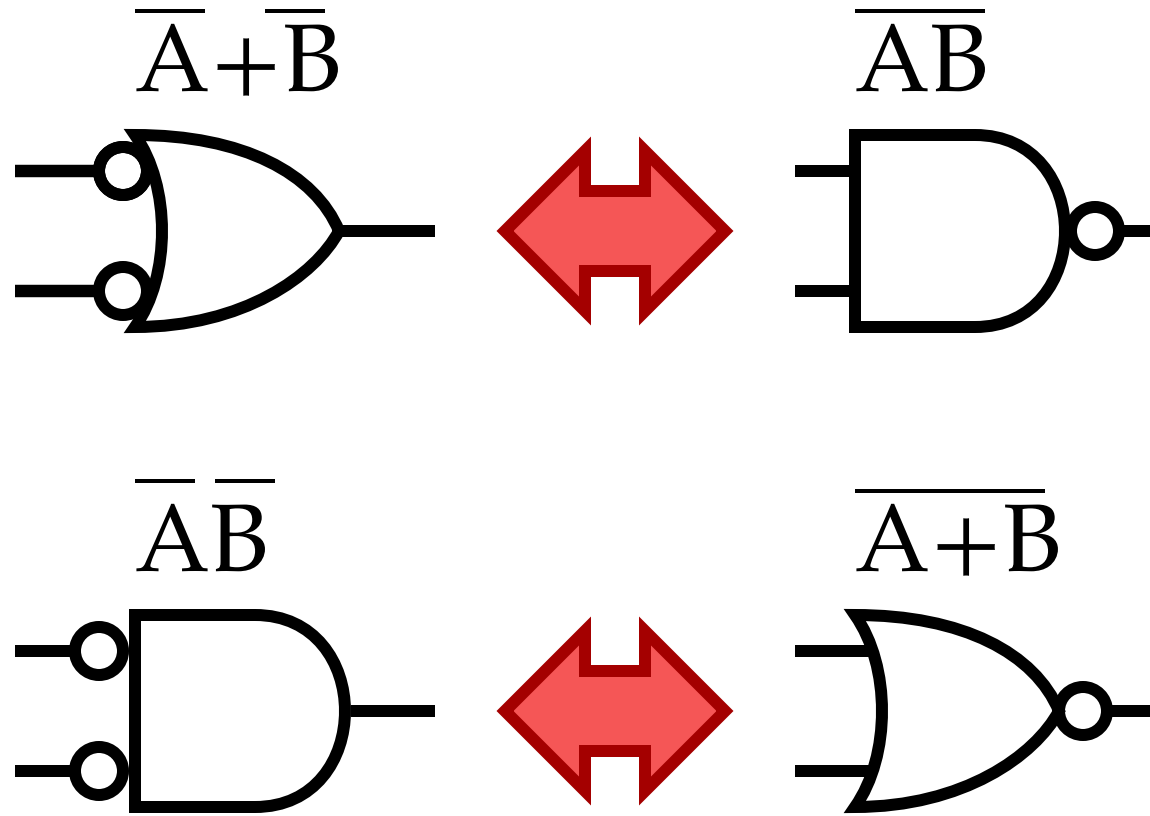
DeMorgan's theorem

$$\overline{A} + \overline{B} = \overline{AB}$$

$$\overline{A} \overline{B} = \overline{A + B}$$

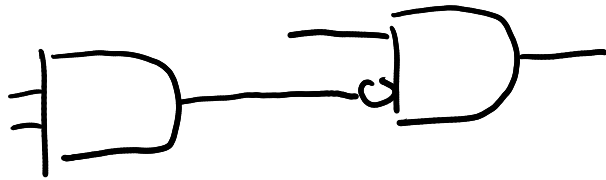
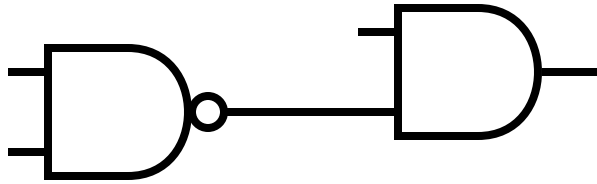
		NAND!	NOR
A	B	$\overline{A} + \overline{B}$	$\overline{A + B}$
0	0	1	1
0	1	1	0
1	0	1	0
1	1	0	0

DeMorgan's theorem on a schematic



Moving bubbles

A bubble just represents inversion, so it can be moved up or down a wire without changing the function.



Involution with bubbles

$$\overline{\overline{A}} = A$$



Bubble pushing

Modify this circuit to use NAND/NOR/NOT instead of AND/OR

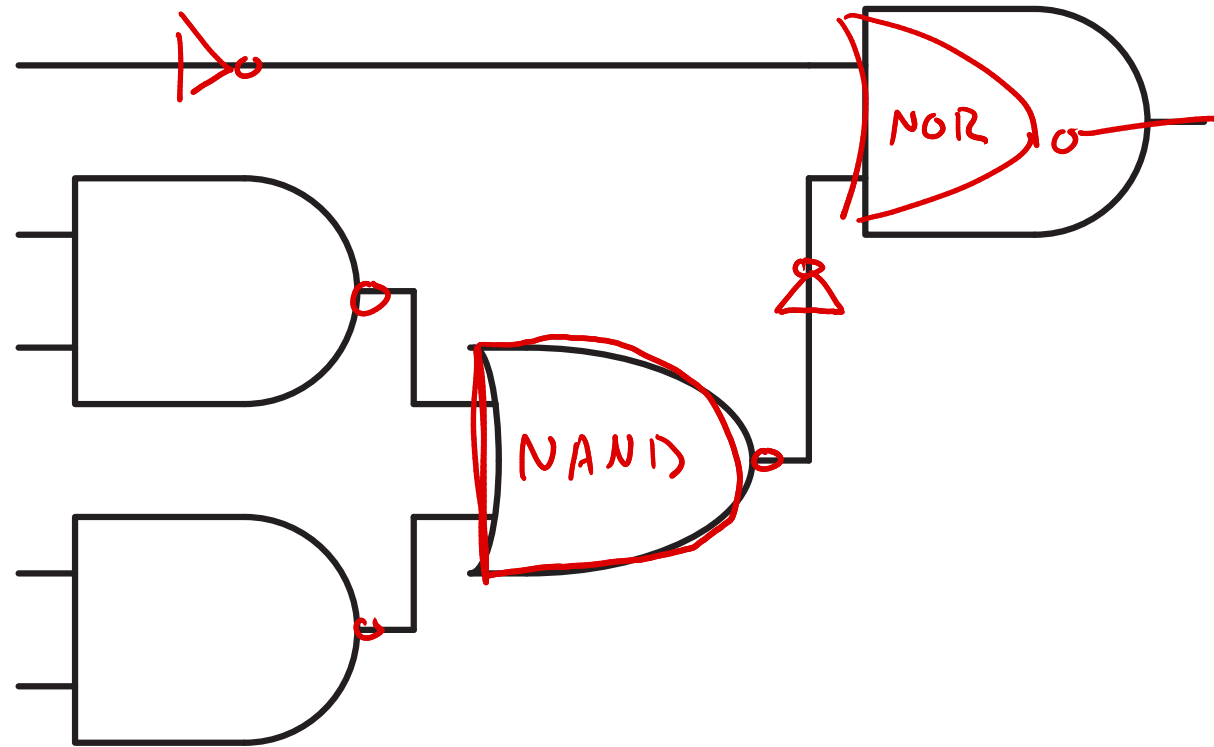


Figure 2.36 from the textbook

Bubble pushing

Modify this circuit to use NAND/NOR/NOT instead of AND/OR

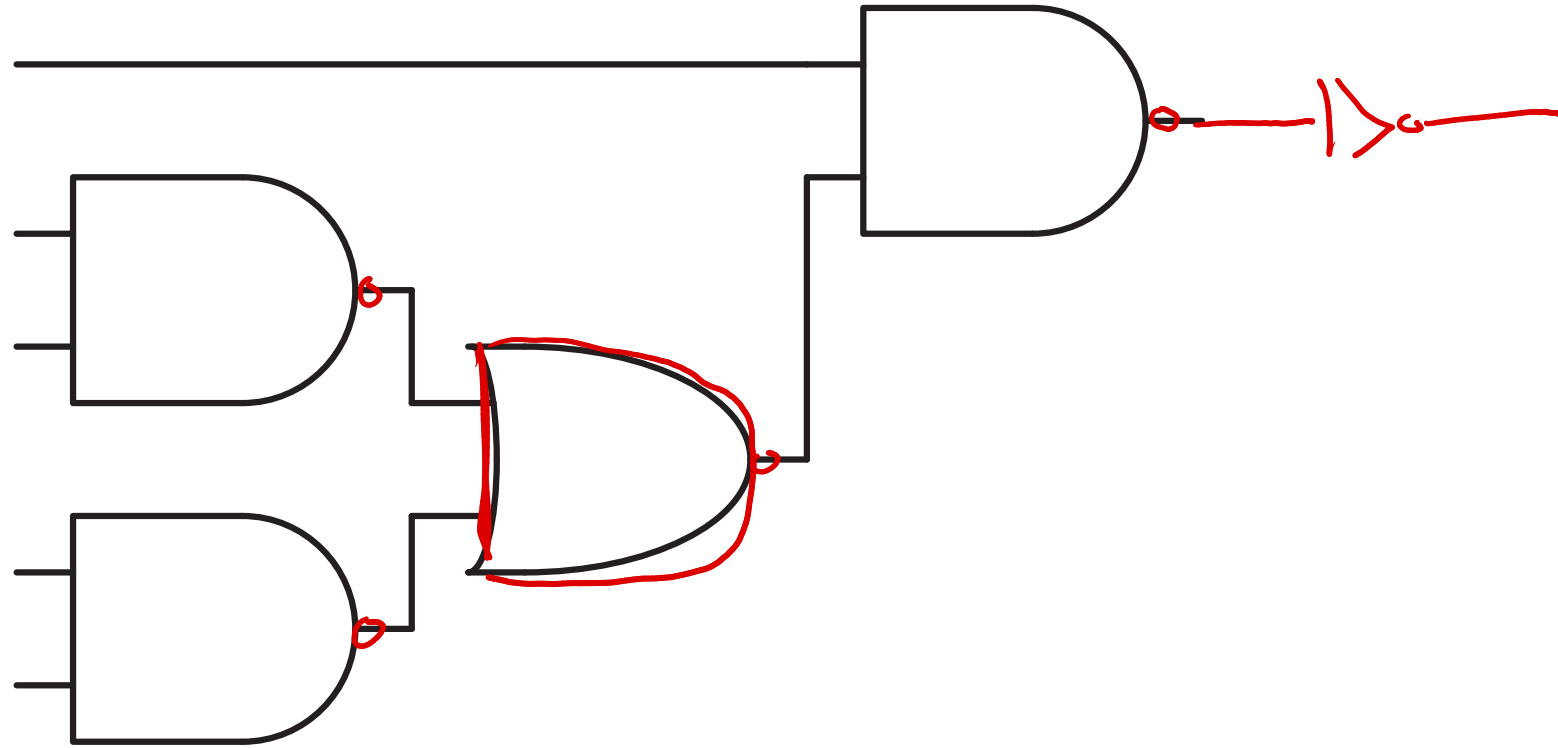
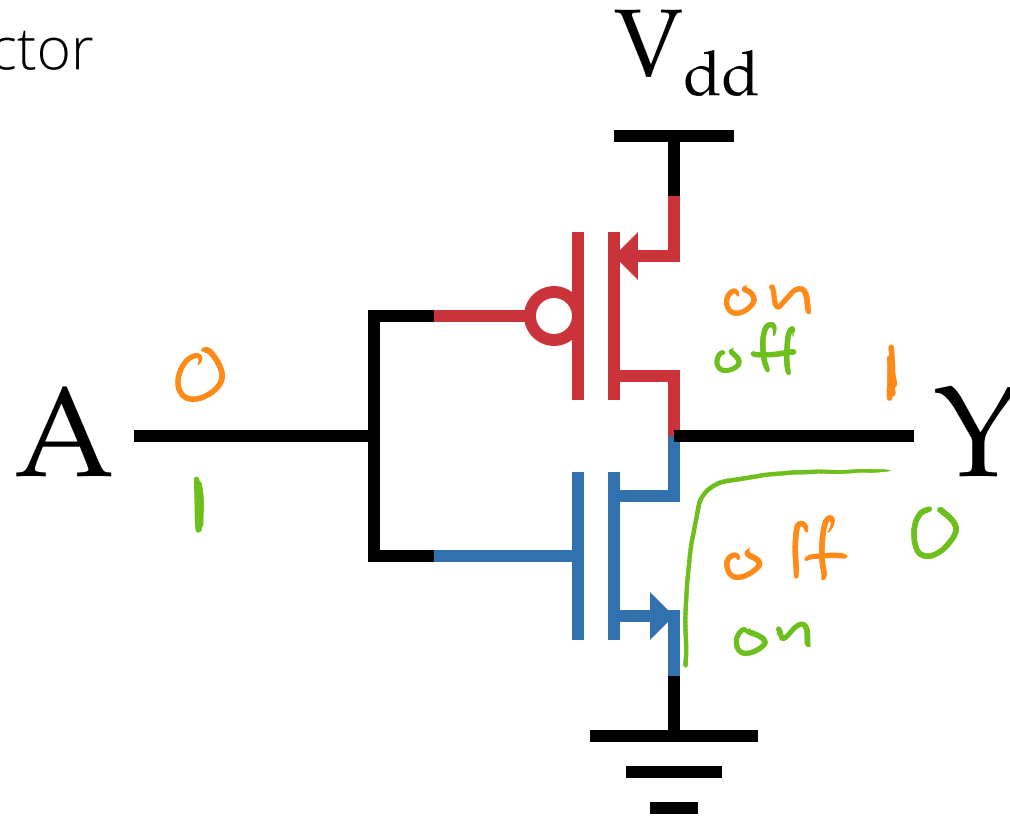


Figure 2.36 from the textbook

A simple gate

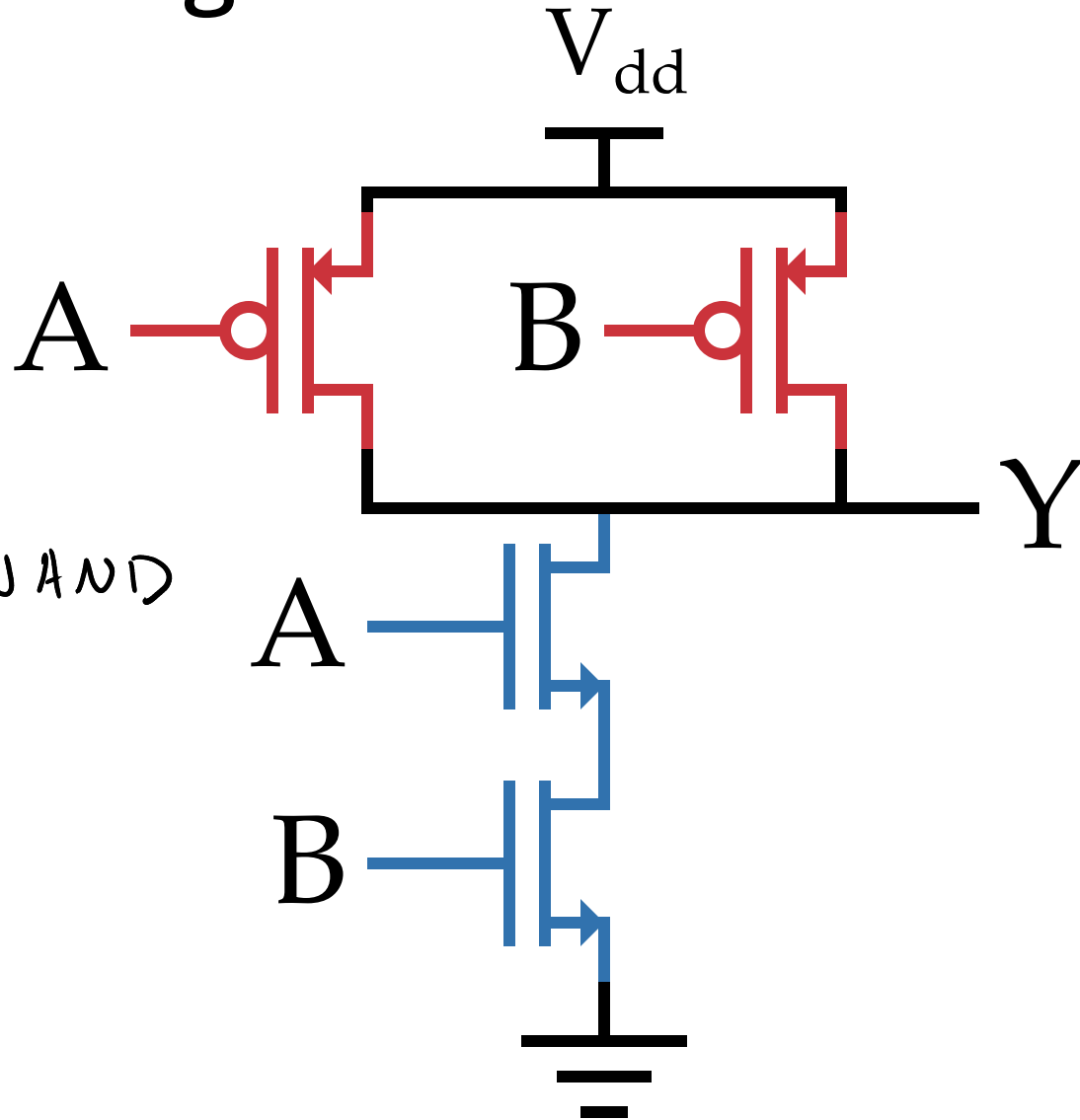
CMOS: *complimentary*
metal-oxide semiconductor



PMOS: "closed" when gate is **low**.

NMOS: "closed" when gate is **high**.

A more complex gate



A	B	Y
0	0	1
0	1	1
1	0	1
1	1	0

NAND

What about the consensus theorem?

It's pretty clever, and you can illustrate it with a Venn diagram but it's not super helpful...

For Thursday

1. Read the book (2.6-2.7) and complete the reading check on Canvas
Reading check is due at **11AM** the day of class, so I can review it
2. Lab sections are posted, plan to come to lab this week
3. Gradescope invites sent; submit your lab notes